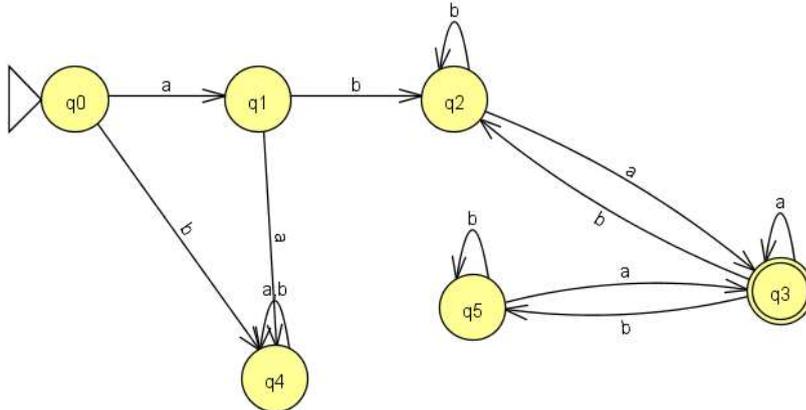


1. Construct a DFA over the language set L for  $\Sigma = \{a,b\}$  such that it accepts strings starting with **ab** and ending with **ba**.

First of all, construct the Language set L

$$L = \{ \mathbf{aba}, \mathbf{abbba}, \mathbf{abababa}, \mathbf{abaabbaaba}, \dots \}$$

Next design the DFA as follows



Construct the Transition Table as per the below formula:

$$\delta(q_0, a) = q_1$$

$$\delta(q_0, b) = q_5 \text{ and so on}$$

$\delta$	a	b
$\rightarrow q_0$	q <sub>1</sub>	q <sub>5</sub>
q <sub>1</sub>	q <sub>4</sub>	q <sub>2</sub>
q <sub>2</sub>	q <sub>3</sub>	q <sub>2</sub>
q <sub>3</sub>	q <sub>3</sub>	q <sub>5</sub>
q <sub>4</sub>	q <sub>4</sub>	q <sub>4</sub>
q <sub>5</sub>	q <sub>3</sub>	q <sub>5</sub>

$$A = \{ Q = \{q_0, q_1, q_2, q_3, q_4, q_5\},$$

$$\Sigma = \{a,b\},$$

$$\delta,$$

$$q_0,$$

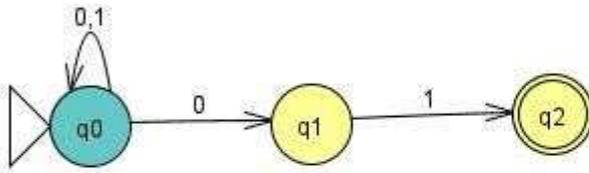
$$F = \{ q_5 \}$$

2. Construct an NFA over the language set L for  $\Sigma = \{0,1\}$  such that it accepts strings ending with **01**.

First of all, construct the Language set L

$$L = \{01, 00001, 111101, 0101001, 10101001, \dots\}$$

Next design the NFA as follows



Construct the Transition Table as per the below formula:

$$\delta(q_0, 0) = q_1$$

$$\delta(q_1, 1) = q_2 \quad \text{and so on}$$

$\delta$	0	1
$\rightarrow q_0$	$q_1$	$\phi$
$q_1$	$\phi$	$q_2$
$q_2$	$\phi$	$\phi$

$$N = \{Q = \{q_0, q_1, q_2\},$$

$$\Sigma = \{0,1\},$$

$$\delta,$$

$$q_0,$$

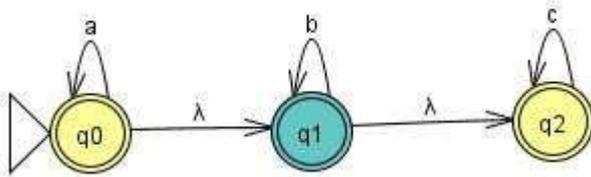
$$F = \{q_2\}$$

3. Construct an  $\epsilon$ -NFA over the language set L for  $\Sigma = \{a,b\}$  such that it accepts strings containing zero or more a's followed by zero or more b's followed by zero or more c's.

Firstly define the Language set L

**L =  $\{\epsilon, a, b, c, aaa, bbb, ccc, abc, aaabbbccc \dots\}$**

Next design the  $\epsilon$  - NFA as follows



Construct the Transition Table as per the below formula:

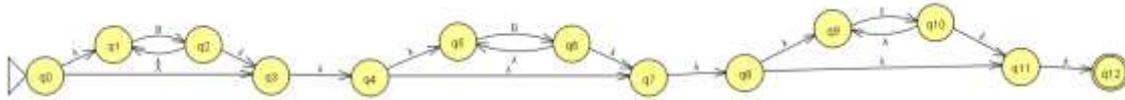
$$\delta(q_0, a) = q_0$$

$$\delta(q_1, b) = q_1 \quad \text{and so on}$$

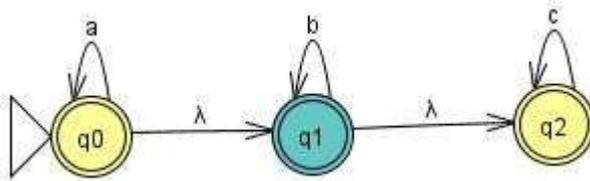
$\delta$	a	b	c	$\epsilon$
$\rightarrow q_0$	$q_0$	$\phi$	$\phi$	$q_1$
$q_1$	$\phi$	$q_1$	$\phi$	$q_2$
$q_2$	$\phi$	$\phi$	$q_2$	$\phi$

$$N = \{\{q_0, q_1, q_2\}, \{a,b\}, \delta, q_0, \{q_0, q_1, q_2\}\}$$

A more formal way of framing  $\epsilon$ -NFA for the above problem



4. Convert the following  $\epsilon$ -NFA into its DFA form



Firstly, draw the transition table of above automata as per the below formula:

$$\delta(q_0, a) = q_0$$

$$\delta(q_1, b) = q_1 \text{ and so on}$$

$\delta$	a	b	c	$\epsilon$
$\rightarrow q_0$	$q_0$	$\phi$	$\phi$	$q_1$
$q_1$	$\phi$	$q_1$	$\phi$	$q_2$
$q_2$	$\phi$	$\phi$	$q_2$	$\phi$

Determine the  $\epsilon$ - closure of each state as follows:

$\epsilon$ - closure of

$$q_0 = \{q_0, q_1, q_2\}$$

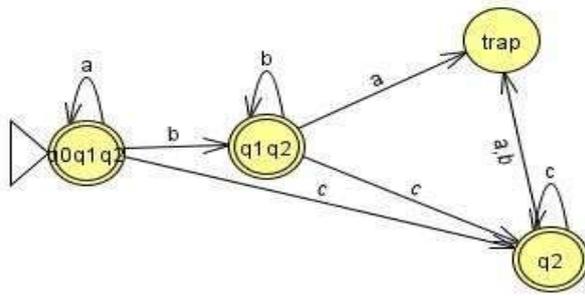
$$q_1 = \{q_1, q_2\}$$

$$q_2 = \{q_2\}$$

Put the transition table for the DFA referring to above  $\epsilon$ - closures.

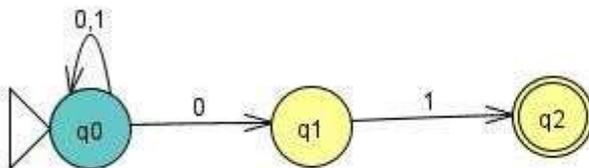
$\delta$	a	b	c
$\rightarrow \{q_0, q_1, q_2\}$	$\{q_0, q_1, q_2\}$	$\{q_1, q_2\}$	$\{q_2\}$
$\{q_1, q_2\}$	$\phi$	$\{q_1, q_2\}$	$\{q_2\}$
$\{q_2\}$	$\phi$	$\phi$	$\{q_2\}$

Now put the final converted DFA as below



$$N = \{ \{ \{q_0, q_1, q_2\} \{q_1, q_2\}, \{q_2\} \}, \{a, b\}, \delta, q_0q_1q_2, \{ \{q_0, q_1, q_2\} \{q_1, q_2\}, \{q_2\} \} \}$$

5. Convert the below shown NFA to its DFA form



Construct the Transition Table as per the below formula:

$$\delta(q_0, 0) = q_1$$

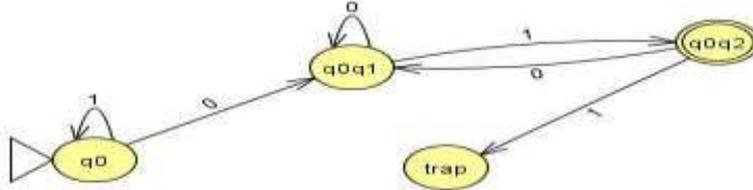
$$\delta(q_1, 1) = q_2 \text{ and so on}$$

$\delta$	0	1
$\rightarrow q_0$	$q_0, q_1$	$q_0$
$q_1$	$\phi$	$q_2$
$q_2$	$\phi$	$\phi$

Construct the transition table for its DFA

$\delta$	0	1
$\rightarrow q_0$	$q_0, q_1$	$q_0$
$q_0, q_1$	$q_0, q_1$	$q_0, q_2$
$q_0, q_2$	$q_0, q_1$	$q_0$

Now put the final DFA as below



Whichever state contains  $q_2$  will be the final state and initial state remains as it is.